

Hot Shots

The first task which must be done before any machine code routines can be loaded on to Basic is to find somewhere to put them. High-memory sites are the usual place for not bytes and on the Amstrad that means looking at locations near to &BFFF, since the top 16K (&C000-FFFF) is occupied by both screen RAM and the Basic ROM. The top of the usable memory, HIMEM as it's called, is somewhere near to high as &BFFF, because Basic needs workspace, the Z-80 processor needs a stack, and the tape-on-tape - flag system likes to take 4K of RAM from time to time.

Also, if you are in the habit of defining dozens of little space invader characters, you could say goodbye to another 1K. To use how low HIMEM can get on your machine, to set the computer and run the following program:

```
100 PRINT HIMEM:GOTO 100
110 GOTO 100
120 PRINT HIMEM:GOTO 100
130 GOTO 100
140 PRINT HIMEM:GOTO 100
150 GOTO 100
```

The final HIMEM figure will be somewhere between &BFFF and &AFFF. The short routines used to illustrate in this article will be placed between &B000 and &B000 and HIMEM should be set to &BFFF, so that the routines being above are protected from Basic. MEMOBY &BFFF is the command needed to do that.

Both the BASIC and the CALL use the same system to deliver information to the installed machine code routine and to receive back information from it. The system is flexible, efficient and even generous; 64 assorted numerical values and string characters can be interchanged by any single BASIC or CALL. The key to its efficiency lies in the use of a special stack set up by Basic and modified by the Z-80 IX register.

To show how it works there is a routine which loads a 32-bit number into two adjacent memory locations, the double

CALL mc,&B000,&B000: Executes 8000 M0000,&B000
16 bytes of code used.

```
OP 7      : 7 parameters are being passed (&B000 & &B000) :
REG 7     : Register A should hold the number of parameters. :P007
REG 04    : If not 1, we cannot continue and must return! :C04
LD 0,110+0 : The first parameter, &B000, is available at 110+0 :B0000
LD 0,110+1 : so get it and put it into 0. :B0001
LD 0,110+0 : The last parameter, &B000, is available at 110+0 :B00000
LD 04,0    : so put the LD into register 0. :C04
LD 04,0    : and then into &B000. :C04
LD 0,110+1 : Now the 655 into register 0. :B0001
INC 0      : Increase 0, by push to &B001. :C01
LD 04,0    : and put 655 into &B000. :C04
RET       : 655 complete, so return to &B000. :C04
```

POKE, sometimes referred to as a DPOKE. The routine - figure one - is shown as a CALL from Basic. It can be changed into an BASIC later.

From this example three rules are apparent. The number of parameters passed is transferred to the A register. The last parameter typed on the Basic line is pointed to by IX+0 and IX+1. Other parameters, if there are any, follow in pairs and in each pair the lower IX, i.e., the "even" IX - holds the Least Significant Byte of information.

Values are passed back to Basic via an integer variable, which must be declared before the CALL is made. Double PEEKs are DPEEKs - what else could they be? - and if a machine code routine is constructed to perform a DPEEK, the value found in memory

must be passed back to Basic. Figure two shows the DPEEK code, as a simple CALL.

The DPEEK code illustrates two more rules. The direction of travel of information is denoted by the @ symbol. Thus, CALL mc,A@, sends the value contained in A@ to the machine code routine, whereas CALL mc,@A@, sends nothing from A@ but accepts a machine code generated value into A@. For an @A@ transfer, the IX pointers indicate indirectly where to place the supplied value, which is assumed to be an unsigned integer.

A string can also be transferred and the format for both sending and receiving is A@="M"CALL mc,@A@. The IX pair involved at the machine code end of the transfer will produce a number

which is the address of those bytes elsewhere in memory, it is those which describe the string. At the first of those bytes is the length of the string and the next two bytes are the address of the string. Once the address of the string is determined, the machine code routine can examine the string, copy it or alter it.

The BASIC system, in its simplest interpretations, is a way of giving names to machine code routines so they do not have to be CALLED by their hexadecimal address. Any name of reasonable length can be used, although it should not be the same as one already used by Basic or by an extension ROM utility. The name is preceded by a bar symbol - to be found on the shifted-@ key - so that it is recognised by the operating system as an BASIC

8000 CALL M000,&B000,&B0: Executes 40-0000-00000
16 bytes of code used.

```
OP 7      : 7 parameters are being passed (&B000 & 0) :
REG 7     : Register A should hold the number of parameters. :P007
REG 04    : If not 2, we cannot continue and must return! :C04
LD 0,110+0 : The first parameter, &B000, is available at 110+0 :B0000
LD 0,110+1 : so get it and put it into 0. :B0001
LD 0,110+0 : The last parameter, 0, has an address which can :B00000
              be retrieved from 110+1. :
LD 0,110+1 : 00 has inside that address, :B000010
              three 1s in bytes are not wanted! See last. :
LD 0,04    : This is a machine-code PEEK. It's the LD. :C04
LD 04,0    : Put 0 into the LD of 04. :C01
INC 04     : Point at 655 of 04. :C01
INC 04     : Advance 04, for next PEEK. :C01
LD 0,04    : Retrieve machine-code PEEK, 655 this time. :B0001
LD 04,0    : Put 0 into 655 of 04. :C01
RET       : 655 complete, so return to &B000. :C04
```

and not taken as a typing error. Setting up a single BASIC command such as DOKE uses 20 bytes, excluding the code of the DOKE routine. Further BASIC commands require an overhead of only a few bytes and they are mostly the name of the routine. Figure three shows the assembly listing for four BASICs.

It is not difficult to think of routines to add to Amstrad Basic and indeed, with the Amstrad Centre Programmer Guide which has the addresses and information about all the ROM sub-routines of interest to machine code programmers, there is every chance evening's worth of coding to be done. The trouble is that there are a host of extremely interesting BASICs already developed and tucked into commercial ROMs which plug into the back of the Amstrad. ROM BASICs do not take precious bytes away from Basic and are always available, so they are useful commands to have.

ENDGAME

The two final examples are PAN α which moves the cursor of the whole screen sideways and PAN β which moves the whole screen sideways — there is a difference. Suggested values for α are 1-80 and for β 1-28. The effects of PAN α last while a program is running, but normally is restored when Basic returns with the READY prompt. PAN β on the other hand causes more permanent lateral shifts and the screen can be returned to normal only by typing PAN β or simply PAN on its own. Note the space between the BSK name and the first comma, that being the required syntax.

TWIST α looks at a character on the screen at row r and column c and rotates it n times through 90 degrees, which is useful for creating vertical writing for labelling the Y axis of graphs. The machine code does not check the legality of r and c and a co-ordinate outside the limits of a screen mode will have unpredictable effects. It also uses the right-hand storage area used for CHR\$(255), so avoid using

	1 The INITIALISATION code.
	2
	3 Where the addresses of all the routines are held.
	4 A 5-byte area of scratch-ROM.
	5 Before the operating system that the ROM exists.
	6 Basic initialisation, as normal.
	7 18 bytes of code as 4*4.
	8
	9 This is 4 bytes of storage.
	10
	11 The COMMANDABLE takes this form.
	12
	13 3 bytes, the address of the ROM's names.
	14 5 bytes in the form of 03 as an where 000 is the address of the first ROM routine.
	15 The jump for ROM routine number 2.
	16 and so on...
	17 The final jump in our example.
	18 bytes
	19
	20 The VARIABLE itself.
	21
	22
	23 4 bytes, the last having bit 7 set high.
	24 4 bytes.
	25 3 bytes.
	26 3 bytes.
	27 1 byte end of table marker.

```

10 REM DEEK, DOKE, PAN & TWIST      120 FOR I=1 TO 32
20 REM AMSTRAD ROMS R. SAMPSON 1985  140 FOR I=1 TO 40NEXT I
30 REM First load the machine-code  150 IF=44 ROMNEXT I
40 MEMORY 32767                    160 FOR N=16 TO 32
50 FOR A=00000 TO 00000             170 FOR I=1 TO 40NEXT I
60 READ B1:FOR A,VAL I*B1+0000NEXT A  180 IF=00 ROM, B, HINEXT I
70 CALL MEMMAP:REM INITIALISE ROMS  190 FOR N=27 TO 36 STEP-1
80 DELETE 32-80                    200 FOR I=1 TO 40NEXT I
90 REM Demonstration of PAN & TWIST  210 IF=00 ROM, B, HINEXT I
100 POKE 1:BSKDFN 12               220 FOR I=1 TO 25
110 FOR I=1 TO 130                  230 FOR I=1 TO 40
120 PRINT "XXXXXXXXXXXXXXXXXXXX"   240 TWIST C,B,HINEXT I
                                     250 NEXT I

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```

900 DATA 01,10,00,21,02,00,0C,01,0C,09,10,00,0C,2F,00,0C
901 DATA 42,00,0C,52,00,0C,70,00,44,40,40,0C,44,45,45,00
902 DATA 50,41,0C,54,57,49,53,04,00,FE,02,00,00,01,02,00
903 DATA 64,02,00,70,00,77,00,70,01,23,77,09,FE,0C,00,00
904 DATA 80,02,00,84,03,70,12,13,23,70,17,09,07,70,23,FE
905 DATA 82,20,04,00,03,00,00,70,FE,01,00,FE,10,00,07,05
906 DATA 0E,FE,01,00,00,3E,02,0C,79,01,00,00,FE,03,79,09
907 DATA 20,70,10,00,FE,0C,00,00,FF,0C,0C,00,00,72,00,00
908 DATA 00,00,0F,0C,00,64,04,20,20,00,0C,02,00,0C,2C,0C
909 DATA 0D,70,00,00,00,00,70,00,01,01,02,0C,00,00,10,FE,FE
910 DATA 0D,70,00,01,02,00,24,2C,0D,70,00,3E,FF,0C,50,00
911 DATA 01,0C,70,00,09,05,00,00,00,FE,0C,00,00,00,02,00,00
912 DATA 00,30,17,27,10,FE,FE,FE,00,FE,02,00,00,00,41,FE
913 DATA 20,77,10,FE,01,FE,0C

```

that character as a user-defined graphic. The listing of figure four loads the code for all four ROMs and then deletes lines 20

to 80, which means that subsequent ROMs do not re-initialise the ROMs and so confuse the computer. The pro-

gram goes on to demonstrate PAN and TWIST: what you do with your ROMs after that is entirely for you.